

Question 1

Let $k=6$. Assume a cache with initial contents 2,3,4,5,6,7 and with at least one item not in the cache. Give an example sequence of length at least 7 of cache requests

1, 2, 3, 4, 5, 6, 7, 1, 2, 3, 4, 5

and a sequence of random choices by the Marker algorithm so that its competitiveness on this sequence with these choices would be greater than H_6 .

Using the sequence above, the Marker algorithm goes as follows:

Cache before request	Marker bits before request	Request	Cache after request	Marker bits after request
234567	000000	1 (miss)	234167	000100
234167	000100	2, 3, 4 (hits)	234167	111100
234167	111100	5 (miss)	234157	111110
234157	111110	6 (miss)	234156	111111
Reset				
234156	000000	7 (miss)	734156	100000
734156	100000	1 (hit)	734156	100100
734156	100100	2 (miss)	734152	100101
734152	100101	3, 4, 5 (hits)	734152	111111

5 misses occurred for this request sequence using Marker algorithm.

The minimum number of misses can be achieved by MIN as follows:

Cache before request	Request	Cache after request
234567	1 (miss)	234561
234561	2, 3, 4, 5, 6 (hits)	234561
234561	7 (miss)	234571
234571	1, 2, 3, 4, 5 (hits)	234571

⇒ 2 misses occurred for the same request sequence for MIN

$$H_6 = 1 / 1 + 1 / 2 + \dots + 1 / 6 = 49 / 20$$

$$C_{\text{Marker}} = 5 / 2 = 50 / 20$$

$$\Rightarrow C_{\text{Marker}} > H_6 (\text{Q.E.D})$$

Question 2.1

Use the extended Euclidean algorithm to find the multiplicative inverse of $14 \text{ mod } 2145$.

```
Extended-Euclid(2145, 14):
    if (14 == 0) then return (14, 1, 0)
    (d', x', y') = Extended-Euclid(14, 3)
    Extended-Euclid(14, 3):
        if (3 == 0) then return (14, 1, 0)
        (d', x', y') = Extended-Euclid(3, 2)
        Extended-Euclid(3, 2):
            if (2 == 0) then return (3, 1, 0)
            (d', x', y') = Extended-Euclid(2, 1)
            Extended-Euclid(2, 1):
                if (1 == 0) then return (2, 1, 0)
                (d', x', y') = Extended-Euclid(1, 0)
                Extended-Euclid(1, 0):
                    if (0 == 0) then return (1, 1, 0)
                    (d', x', y') = (1, 1, 0)
                    (d, x, y) = (1, 0, 1 - 2 * 0)
                    return (1, 0, 1)
                (d', x', y') = (1, 0, 1)
                (d, x, y) = (1, 1, 0 - 1 * 1)
                return (1, 1, -1)
            (d', x', y') = (1, 1, -1)
            (d, x, y) = (1, -1, 1 - 4 * -1)
            return (1, -1, 5)
        (d', x', y') = (1, -1, 5)
        (d, x, y) = (1, 5, -1 - 153 * 5)
        return (1, 5, -766)
```

Therefore, the multiplicative inverse for $14 \text{ mod } 2145$ is -766 (or $1379 \text{ mod } 2145$)

Question 2.2

Solve $6x \equiv 9 \pmod{33}$ for all solutions.

For equation $6x \equiv 9 \pmod{33}$, there are either $\gcd(6, 33) = 3$ solutions or none.

```
Modular-Linear-Equation-Solver(6, 9, 33):
    (d, x', y') = Extended-Euclid(6, 33) # see calculation at \(\*\)
    (d, x', y') = (3, -5, 1)
    solution = []
    if (3 | 9):
        x_0 = -5 * (9/3) mod 33
        x_0 = 18
        for i in [0, 3):
            solution.append(18 + i * (33/3)) mod n
            i = 0:
                solution = [18]
            i = 1:
                solution = [18, 29]
            i = 2:
                solution = [18, 29, (18+22) mod 33]
                = [18, 29, 7]

    return solution
```

Solution for $6x \equiv 9 \pmod{33}$ is [7, 18, 29]

(*) Euclid(6, 33) calculation

```
Extended-Euclid(6, 33):
    if(33 == 0) return (6, 1, 0)
    (d', x', y') = Extended-Euclid(33, 6)
    Extended-Euclid(33, 6):
        if(6 == 0) return (33, 1, 0)
        (d', x', y') = Extended-Euclid(6, 3)
    Extended-Euclid(6, 3):
        if (3 == 0) return (6, 1, 0)
        (d', x', y') = Extended-Euclid(3, 0)
        Extended-Euclid(3, 0):
            if (0 == 0) return (3, 1, 0)
            (d', x', y') = (3, 1, 0)
            (d, x, y) = (3, 0, 1 - 2 * 0)
            return (3, 0, 1)
        (d', x', y') = (3, 0, 1)
        (d, x, y) = (3, 1, 0 - 5 * 1)
        return (3, 1, -5)
    (d', x', y') = (3, 1, -5)
    (d, x, y) = (3, -5, 1 - 0 * -5)
    return (3, -5, 1)
```

Question 3

Using the Chinese Remainder theorem, determine a number $x \bmod 2145$ that satisfies $x \equiv 2 \bmod 3$, $x \equiv 3 \bmod 5$, and $x \equiv 4 \bmod 11$, and $x \equiv 5 \bmod 13$

$$\begin{aligned} n_i &\Rightarrow [3, 5, 11, 13] \\ m_i &\Rightarrow [2145/3, 2145/5, 2145/11, 2145/13] \\ &= [715, 429, 195, 165] \end{aligned}$$

Run Extended Euclid (m_i, n_i) (denoted as EE)

i = 1

EE(715, 3)

$$\begin{aligned} &\Rightarrow \text{EE}(3, 1) \\ &\Rightarrow \text{EE}(1, 0) \\ &\Rightarrow (1, 1, 0) \\ &\Rightarrow (1, 0, 1 - 3 * 0) \\ &\Rightarrow (1, 0, 1) \\ &\Rightarrow (1, 1, 0 - 238 * 1) \\ &\Rightarrow (1, 1, -238) \end{aligned}$$

$$t_1 = 715^{-1} = 1 \bmod 3$$

$$c_1 = m_1 * t_1 = 715 * 1 = 715$$

i = 2

EE(429, 5)

$$\begin{aligned} &\Rightarrow \text{EE}(5, 4) \\ &\Rightarrow \text{EE}(4, 1) \\ &\Rightarrow \text{EE}(1, 0) \\ &\Rightarrow (1, 1, 0) \\ &\Rightarrow (1, 0, 1 - 4 * 0) \\ &\Rightarrow (1, 0, 1) \\ &\Rightarrow (1, 1, 0 - 5 // 4 * 1) \\ &\Rightarrow (1, 1, -1) \\ &\Rightarrow (1, -1, 1 - 429//5 * -1) \\ &\Rightarrow (1, -1, 86) \end{aligned}$$

$$t_2 = 429^{-1} = -1 \bmod 5$$

$$c_2 = 429 * -1 = -429$$

i = 3 (EE was called similar to above cases, simplified to the inline format as follows)

$$\text{EE}(195, 11) \Rightarrow \text{EE}(11, 8) \Rightarrow \text{EE}(8, 3) \Rightarrow \text{EE}(3, 2) \Rightarrow \text{EE}(2, 1) \Rightarrow \text{EE}(1, 0) \curvearrowright$$

$$(1, -4, 71) \leq (1, 3, -4) \leq (1, -1, 3) \leq (1, 1, -1) \leq (1, 0, 1) \leq (1, 1, 0) \curvearrowright$$

$$t_3 = 195^{-1} = -4 \bmod 11$$

$$c_3 = 195 * -4 = -780$$

i=4

$$\begin{aligned} \text{EE}(165, 13) &\Rightarrow \text{EE}(13, 9) \Rightarrow \text{EE}(9, 4) \Rightarrow \text{EE}(4, 1) \Rightarrow \text{EE}(1, 0) \rightarrow \\ (1, 3, -38) &\leq (1, -2, 3) \leq (1, 1, -2) \leq (1, 0, -1) \leq (1, 1, 0) \end{aligned}$$

$t_4 = 165^{-1} = 3 \bmod 13$
 $c_4 = 165 * 3 = 495$

$$\begin{aligned} a_i &\Rightarrow [2, 3, 4, 5] \\ c_i &\Rightarrow [715, -429, -780, 495] \end{aligned}$$

$$\begin{aligned} x &= 2 * 715 + 3 * -429 + 4 * -780 + 5 * 495 \\ &= 1430 - 1287 - 3120 + 2475 \\ &= -502 \bmod 2145 \\ &= \underline{\underline{1643 \bmod 2145}} \end{aligned}$$

Check:

$$\begin{aligned} 1643 \div 3 &= 547 \dots 2 \\ 1643 \div 5 &= 328 \dots 3 \\ 1643 \div 11 &= 149 \dots 4 \\ 1643 \div 13 &= 126 \dots 5 \end{aligned}$$

Question 4

Suppose $p=11, q=17$. If we choose $e=3$, what would be the RSA public and private keys?

$$p = 11, q = 17, e = 3$$

$$n = pq = 11 * 17 = 187$$

$\phi(n) = (p-1)(q-1) = 160$ was indeed relatively prime to 3

$$\text{EE}(160, 3) \Rightarrow \text{EE}(3, 1) \Rightarrow \text{EE}(1, 0)$$

(1, 1, -53) <= (1, 0, 1) <= (1, 1, 0)

$$160 * 1 + 3 * -53 = 1$$

$$3 * -53 = 1 \bmod 160, d = -53 \bmod 160 = \mathbf{107 \bmod 160}$$

$$\text{pubkey} = (3, 187)$$

$$\text{secret} = (107, 187)$$

$$M = 83$$

Show the result of encrypting with the private key, the message 83. Show the steps in decrypting it, to get the original number back.

Encrypt

Cipher_text = Message^e mod n
 $C = 83^{107} \bmod 187 = 162$ (Using the *expMod* function in my Java program)

```
1 static int expMod(int x, int y, int p){  
2     if (x == 0)  
3         return 0;  
4     if (y == 0)  
5         return 1;  
6  
7     long tmp;  
8     if (y % 2 == 0){  
9         tmp = expMod(x, y / 2, p);  
10        tmp = (tmp * tmp) % p;  
11    }else{  
12        tmp = x % p;  
13        tmp = (tmp * expMod(x, y - 1, p) % p) % p;  
14    }  
15  
16    return (int)((tmp + p) % p);  
17}
```

Decrypt

Plain_text = Cipher_text^d mod n
 $P = 162^3 \bmod 187 = \mathbf{83}$ (equals to original message!)

Coding question

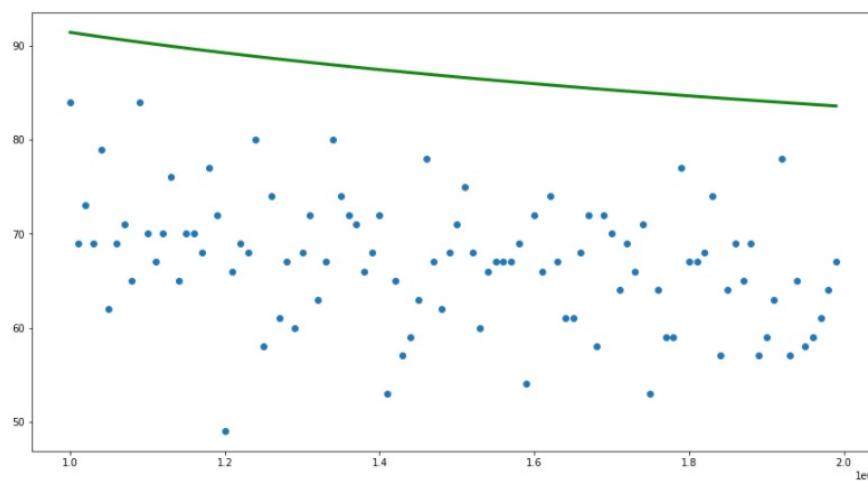
According to Brun's argument¹, the number of twin primes grow in $O(N/\log^2(N))$, so I did linear regression on $N/\log^2(N)$. In this experiment, the number of primes could be bounded by $2.04 \times N/\log^2(N)$ as follows.

```
In [33]: df = pd.read_csv('result10.csv')
x = list(zip(df.nDividedByLogSquared, [0 for _ in range(100)]))
y = df[df.predict == 0].prime
reg = LinearRegression().fit(x, y)
print(reg.coef_)
df.predict = df.nDividedByLogSquared * reg.coef_[0]

[2.03920101 0.         ]

In [34]: plt.figure(figsize = (15, 8))
plt.scatter(x = df.bins, y = df.prime)
plt.plot(df.bins, df.predict, linewidth = 3, c = 'green')

Out[34]: <matplotlib.lines.Line2D at 0x7fd60b8f22b0>
```



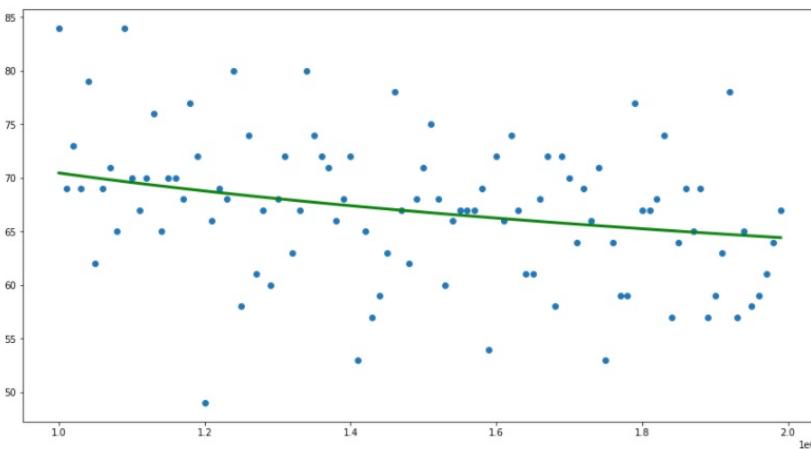
If we were to predict the number of primes in a range, $1.57 \times N/\log^2(N)$ could be a better fit. (the data was too scattered for a good function without large bias)

```
In [35]: df = pd.read_csv('result10.csv')
x = list(zip(df.nDividedByLogSquared, [0 for _ in range(100)]))
y = df[df.predict == 0].prime
reg = LinearRegression(fit_intercept = False).fit(x, y)
print(reg.coef_)
df.predict = df.nDividedByLogSquared * reg.coef_[0]

[1.57141764 0.         ]

In [36]: plt.figure(figsize = (15, 8))
plt.scatter(x = df.bins, y = df.prime)
plt.plot(df.bins, df.predict, linewidth = 3, c = 'green')

Out[36]: <matplotlib.lines.Line2D at 0x7fd60baeb640>
```



¹ https://en.wikipedia.org/wiki/Brun%27s_theorem