1st Order Logic

Giving meaning to 1st order formulas:

Need to have a set M called the universe that our variables range over. For each constant c in our language, we need to give a value C^m in M. For each function f in our language need to give an actual function determined on f^m

```
For each predicate symbol P
P^{m}:M \rightarrow \{T, \bot\}
```

All of the above information together is called a structure/model. Usually denoted by M.

Can almost now talk about the meaning of a formula F in a given structure M.

Last thing, in F(x1, ..., xn): (Exists x)G(x, y) x is a bound variable, y is unbound

A variable/object assignment v is a map from unbound variables to elements of M.

Write $M \models F[v]$ (*M* entails F with assignment v)

What this means...

If F is an atomic formula calculate the value in M of each of the terms in F. Plug these values into the predicate P^m and see if outputs T or F (\bot)

If F := NOT G then F is true in M, v iff not $M \models G[v]$

If F := G OR H Then $M \models F[v]$ iff $M \models G[v]$ or $M \models H[v]$

If F := G AND H then $M \models F[v]$ iff $M \models G[v]$ and $M \models H[v]$

If F := (Exists x) G then $M \models F[v]$ iff there is some way to map $x \rightarrow x^m$ in M such that v is extended by this additional mapping give $M \models G [x \rightarrow x^m, v]$

If F := (Forall x) G then M |= F[v] iff for all ways to map $x \rightarrow x^m$ in M M |= G [x->x^m, v]

Write M \models F if for every variable assignment, v, M \models F[v]

In AI had a notion of a knowledge base entailing a formula Write KB |= F to mean for all structures M, such that M |= KB we also have M |= F

Proofs in 1st Order Logic

Very much like proofs in propositional logic except now we have some additional axioms.

Example additional axioms

(Forall x) (NOT P) => NOT((Exists x) P) (for every pig, it can't fly = not there is a pig that can fly) NOT(Forall x P) =>(Exists x) NOT P (Not for every horse it is blue = there is a horse that is not blue)

 $A(t) \Rightarrow (exists x) A(x)$ $A(x) \Rightarrow (forall y) A(y)$

 $(A AND B) \Rightarrow A$

etc.

Example 1st Order Knowledge bases

- a) Any relational database.
- b) Set theory
 Has one constant 0 (empty)
 and predicate symbols C (element of), =
 No function symbols

Knowledge base for set theory consists of formulas like NOT($0 \in 0$) Exists x NOT(x = 0)