

Enigma



- Developed and patented (in 1918) by Arthur Scherbius
- Many variations on basic design
- Eventually adopted by Germany
 - For both military and diplomatic use
 - Many variations used
- Broken by Polish cryptanalysts, late 1930s
- Exploited throughout WWII
 - o By Poles, British, Americans



- Turing was one of Enigma cryptanalysts
- Intelligence from Enigma vital in many battles
 - o D-day disinformation
 - German submarine "wolfpacks"
 - Many other examples
- May have shortened WWII by a year or more
- Germans never realized Enigma broken Why?
 - o British were cautious in use of intelligence
 - But Americans were less so (e.g., submarines)
 - Nazi system discouraged critical analysis...

Enigma

 To encrypt

 Press plaintext letter, ciphertext lights up

 To decrypt

 Press ciphertext letter, plaintext lights up

Electo-mechanical



Enigma Crypto Features

3 rotors

• Set initial positions

Moveable ring on rotor

o Odometer effect

Stecker (plugboard)

• Connect pairs of letters

Reflector

• Static "rotor"



Enigma

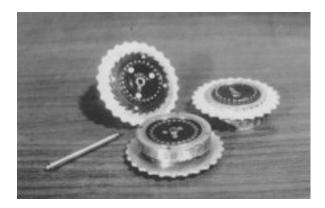
Substitution Cipher

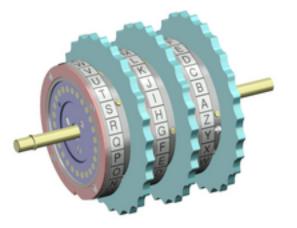
- Enigma is a substitution cipher
- But not a simple substitution
 - Perm changes with each letter typed
- Another name for simple substitution is mono-alphabetic substitution
- Enigma is an example of a poly-alphabetic substitution
- How are Enigma "alphabets" generated?

Enigma Components

Each rotor implements a permutation The reflector is also a permutation o Functions like stecker with 13 cables Rotors operate almost like odometer • Reflector does not rotate • Middle rotor occasionally "double steps" Stecker can have 0 to 13 cables

Enigma Rotors



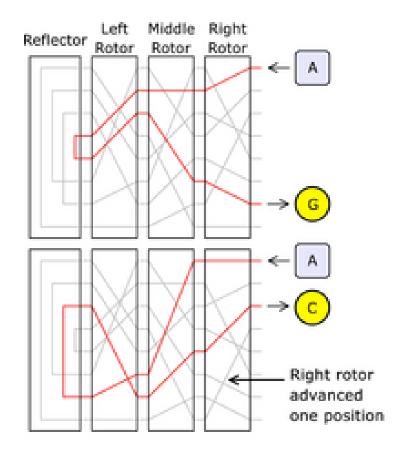


Three rotors

Assembled rotors

Rotors and Reflector

 Each rotor/reflector is a permutation
 Overall effect is a permutation
 Due to odometer effect, overall permutation changes at each step



Why Rotors?

Inverse permutation is easy

• Need inverse perms to decrypt!

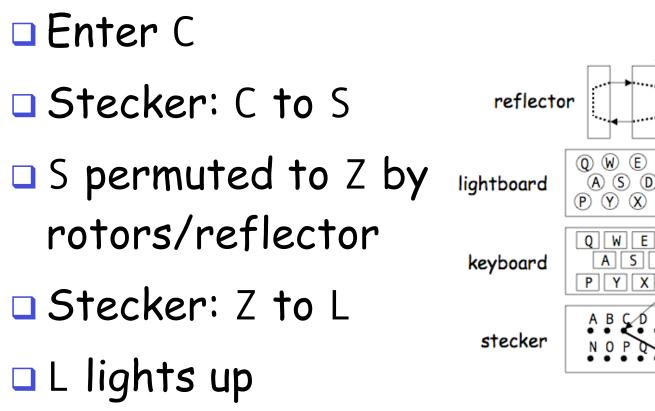
• Pass current thru rotor in opposite direction

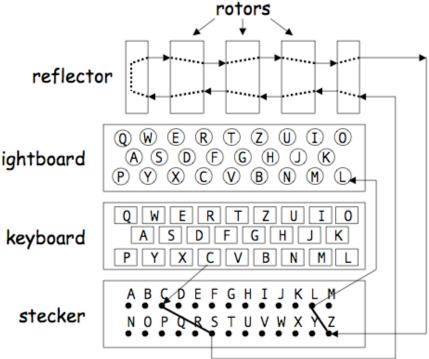
Can decrypt with same machine

• Maybe even with the same settings...

- Rotors provide easy way to generate large number of permutations mechanically
- Otherwise, each perm would have to be wired separately (as in Purple cipher...)

Wiring Diagram





Enigma is Its Own Inverse!

- Suppose at step i, press X and Y lights up
 - Let A = permutation thru reflector
 - Let B = thru leftmost rotor from right to left
 - Let C = thru middle rotor, right to left
 - o Let D = thru rightmost rotor, right to left
- $\Box \text{ Then } Y = S^{-1}D^{-1}C^{-1}B^{-1}ABCDS(X)$
- Where "inverse" is thru the rotor from left to right (inverse permutation)
- Note: reflector is its own inverse

• Only one way to go thru reflector

Inverse Enigma

 Suppose at step i, we have Y = S⁻¹D⁻¹C⁻¹B⁻¹ABCDS(X)
 Then at step i X = S⁻¹D⁻¹C⁻¹B⁻¹ABCDS(Y)
 Since A = A⁻¹

Why is this useful?

Enigma Key?

- What is the Enigma key? Machine settings What can be set? o Choice of rotors Initial position of rotors • Position of movable ring on rotor o Choice of reflector Number of stecker cables
 - Plugging of stecker cables

Enigma Keyspace

Choose rotors

 $\circ 26! \cdot 26! \cdot 26! = 2^{265}$

- Set moveable ring on right 2 rotors
 26 · 26 = 2^{9.4}
- □ Initial position of each rotor o $26 \cdot 26 \cdot 26 = 2^{14.1}$
- Number of cables and plugging of stecker
 Next slide
- Choose of reflector
 - Like stecker with 13 cables...
 - ...since no letter can map to itself

Enigma Key Size

Let F(p) be ways to plug p cables in stecker

- Select 2p of the 26 letters
- Plug first cable into one of these letters
- Then 2p 1 places to plug other end of 1st cable
- Plug in second cable to one of remaining
- Then 2p 3 places to plug other end

o And so on...

□ $F(p) = binomial(26,2p) \cdot (2p-1) \cdot (2p-3) \cdot \cdots \cdot 1$

Enigma Keys: Stecker

F(0) :	= 1
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- F(2) = 44850
- F(4) = 164038875
- F(6) = 100391791500
- F(8) = 10767019638375
- F(10) = 150738274937250

- F(1) = 325
- F(3) = 3453450
- F(5) = 5019589575
- F(7) = 1305093289500
- F(9) = 53835098191875
- F(11) = 205552193096250
- F(12) = 102776096548125 F(13) = 7905853580625
- $F(0) + F(1) + ... + F(13) = 532985208200576 = 2^{48.9}$
- Note that maximum is with 11 cables
- □ Note also that $F(10) = 2^{47.1}$ and $F(13) = 2^{42.8}$

Enigma

Enigma Keys

- Multiply to find total Enigma keys 2²⁶⁵ · 2^{9.4} · 2^{14.1} · 2^{48.9} · 2^{42.8} = 2³⁸⁰
- "Extra" factor of $2^{14.1}$ $2^{265} \cdot 2^{9.4} \cdot 2^{48.9} \cdot 2^{42.8} = 2^{366}$
- Equivalent to a 366 bit key!
- Less than 10⁸⁰ = 2²⁶⁶ atoms in observable universe!
- Unbreakable? Exhaustive key search is certainly out of the question...

In the Real World (ca 1940)

- **5** known rotors: $5 \cdot 4 \cdot 3 = 2^{5.9}$
- □ Moveable rings on 2 rotors: 2^{9.4}
- □ Initial position of 3 rotors: 2^{14.1}
- □ Stecker usually used 10 cables: 247.1
- Only 1 reflector, which was known: 20
- Number of keys "only" about
 25.9 · 29.4 · 214.1 · 247.1 · 20 = 276.5

In the Real World (ca 1940)

Only about 2^{76.5} Enigma keys in practice

Still an astronomical number

• Especially for 1940s technology

But, most of keyspace is due to stecker

□ If we ignore stecker...

• Then only about 2²⁹ keys

• This is small enough to try them all

Attack we discuss "bypasses" stecker

Many different Enigma attacks

• Most depend on German practices...

o ...rather than inherent flaws in Enigma

Original Polish attack is noteworthy

- Some say this is greatest crypto success of war
- Did not know rotors or reflector
- Were able to recover these
- Needed a little bit of espionage...

The attack we discuss here

- o Assumes rotors are known
- Shows flaw in Enigma
- Requires some known plaintext (a "crib" in WWII terminology)
- Practical today, but not quite in WWII

Suppose we have known plaintext (crib) below

 0
 1
 2
 3
 4
 5
 6
 7
 8
 9
 10
 11
 12
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 14
 15
 16
 17
 18
 19
 20
 21
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 19
 20
 21
 22
 23

 Plaintext
 O
 B
 E
 R
 K
 M
 A
 N
 D
 O
 D
 E
 R
 W
 A
 N
 T
 S
 W
 U
 I
 N
 Z

 Ciphertext
 Z

- Let P_i be permutation (except stecker) at step i
- □ S is stecker
 - $O M = S^{-1} P_8 S(A) \implies S(M) = P_8 S(A)$
 - o $E = S^{-1} P_6 S(M) \Rightarrow S(E) = P_6 S(M)$
 - $o A = S^{-1} P_{13} S(E) \Rightarrow S(A) = P_{13} S(E)$
- □ Combine to get "cycle" $P_6P_8P_{13}S(E) = S(E)$

 0
 1
 2
 3
 4
 5
 6
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 12
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 17
 18
 19
 20
 21
 22
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 23
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 20
 21
 22
 23

 Plaintext
 O
 B
 E
 R
 M
 A
 C

Also find the cycle E = S⁻¹ P₃S(R) ⇒ S(E) = P₃S(R) W = S⁻¹ P₁₄S(R) ⇒ S(W) = P₁₄S(R) W = S⁻¹ P₇S(M) ⇒ S(W) = P₇S(M) E = S⁻¹ P₆S(M) ⇒ S(E) = P₆S(M)

• Combine to get $P_6 P_{14}^{-1} P_7 P_6^{-1} S(E) = S(E)$

□ Guess one of 2²⁹ settings of rotors o Then all putative perms P_i are known □ If guess is correct cycles for S(E) hold o If incorrect, only 1/26 chance a cycle holds But we don't know S(E) • So we guess S(E) For correct rotor settings and S(E), • All cycles for S(E) must hold true

- Using only one cycle in S(E), must make 26 guesses and each has 1/26 chance of a match
 - On average, 1 match, for 26 guesses of S(E)
 - Number of "surviving" rotor settings is about 229
- But, if 2 equations for S(E), then 26 guesses for S(E) and only 1/26² chance both cycles hold
 - Reduce possible rotor settings by a factor of 26
 - With enough cycles, will have only 1 rotor setting!
 - In the process, stecker (partially) recovered!
- Divide and conquer!

- Enigma was ahead of it's time
- Weak, largely due to combination of "arbitrary" design features
 - For example, right rotor is "fast" rotor
 - If left rotor is "fast", it's stronger
- Some Enigma variants used by Germans are much harder to attack
 - Variable reflector, stecker, etc.

- Germans confused "physical security" and "statistical security" of cipher
 - Modern ciphers: statistical security is paramount
 - Embodied in Kerckhoffs Principle
- Pre-WWII ciphers, such as codebooks
 - Security depends on codebook remaining secret
 - That is, physical security is everything
- Germans underestimated statistical attacks

Aside...

- Germans had some cryptanalytic success
 - o Often betrayed by Enigma decrypts
- □ In one case, **before** US entry in war
 - British decrypted Enigma message
 - o German's had broken a US diplomatic cipher
 - British tried to convince US not to use the cipher
 - But didn't want to tell Americans about Enigma!

Pre-computers used to attack Enigma

- Most famous, were the
 - o Polish "bomba", British "bombe"
 - Electro-mechanical devices
- British bombe, essentially a bunch of Enigma machines wired together
- Could test lots of keys quickly
- Noisy, prone to break, lots of manual labor