

ALICE SENDS A MESSAGE TO BOB  
SAYING TO MEET HER SOMEWHERE.

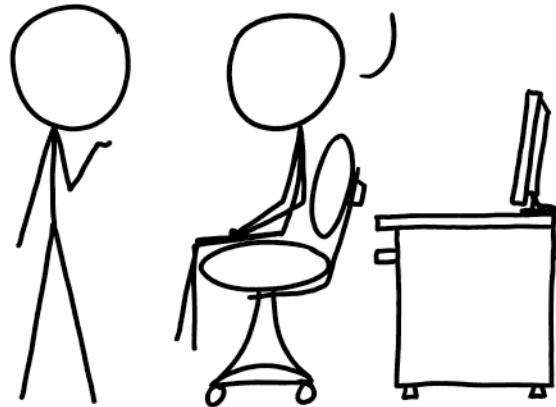
UH HUH.

BUT EVE SEES IT, TOO,  
AND GOES TO THE PLACE.

WITH YOU SO FAR.

BOB IS DELAYED, AND  
ALICE AND EVE MEET.

YEAH?



I'VE DISCOVERED A WAY TO GET COMPUTER  
SCIENTISTS TO LISTEN TO ANY BORING STORY.

<https://xkcd.com/1323/>

# *Cryptocurrencies & Security on the Blockchain*



## **Digicash, Part 1: Blinded Signatures**

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# Lab 3 Review

# DigiCash

- Created by David Chaum
- A *centralized* cryptocurrency
  - Relies on *trusted third party* (TTP)
- Anonymous transactions
- Uses *blind signatures*



# Review: Digital Signatures

- Uses two separate keys
  - Public key is known by everyone
  - Private key known only to the owner
- Encryption
  - Public key encrypts
  - Private key decrypts
- Signing
  - Private key encrypts (signs)
  - Public key decrypts (verifies)

**Public key** is  $(N,e)$

**Private key** is  $d$

To encrypt message  $M$ :

$$C = M^e \bmod N$$

To decrypt ciphertext  $C$ :

$$M = C^d \bmod N$$

**Public key** is  $(N,e)$

**Private key** is  $d$

To sign message  $M$ :

$$S = M^d \bmod N$$

To verify sig.  $S$  for message  $M$ :

$$M' = S^e \bmod N$$

and verify that  $M = M'$

# Blinded signatures

- Motivation: Sender wants a notary to sign a document without revealing the contents.
- Analogy: Signing a piece of carbon paper through an envelope.
- Can the sender get the notary to sign anything it wants?



# Blind Signature Properties

- Signature function and multiplication must be commutative
- Signature remains valid after unblinding
- Signer cannot determine what was signed, until it is unblinded
  - Signer does not even know when a particular document was signed

**Public key** is  $(N, e)$

**Private key** is  $d$

To sign message  $M$ :

$$S = M^d \bmod N$$

To verify sig.  $S$  for message  $M$ :

$$M' = S^e \bmod N$$

and verify that  $M = M'$

# RSA Blind Signature process

- User chooses *blinding factor* B
  - B must be cryptographic quality random number
- User calculates:  
$$M' = M * B^e \bmod N$$
- Bank signs M'  
$$S' = M'^d \bmod N$$
- User removes blinding factor  
$$S = S' / B$$
- Note that  $S' / B = M^d \bmod N$

Why does it work?

$$M' = M * B^e \bmod N$$

$$S' = M'^d \bmod N$$

$$= (M * B^e)^d \bmod N$$

$$= M^d * (B^e)^d \bmod N$$

$$= M^d * B \bmod N$$

$$S' / B = M^d \bmod N$$

# Modular Math

- "Clock math"
- Addition – easy
- Subtraction – easy
- Multiplication – easy
- Division... not so much

# Modular Division

To find  $a / b \bmod N$ :

1. Find the *modular multiplicative inverse* of  $b$

- Denoted  $b^{-1}$
- Might not be defined

2. Return  $a * b^{-1} \bmod N$

# Modular Multiplicative Inverse

- Multiplicative inverse of  $b$  in mod  $N$  math  
Find number  $b^{-1}$  such that
  - $b * b^{-1} = 1 \text{ mod } N$
- If  $b^{-1}$  exists, then the greatest common divisor of  $b$  and  $N$  is 1.
  - $\text{gcd}(b,n) = 1$
- Extended Euclidean algorithm calculates  $b^{-1}$ 
  - Modified version of finding gcd
  - [https://en.wikipedia.org/wiki/Extended\\_Euclidean\\_algorithm](https://en.wikipedia.org/wiki/Extended_Euclidean_algorithm)

# Lab, Part 1: Implement Blind RSA

Download `rsa.js` from the course website. It demonstrates the math for the RSA algorithm.

Add `blind` and `unblind` functions for working with signatures. Test out your solution, then paste these functions into Canvas.

For unblinding, you will need to do modular division. You may find `modularDivision.js` useful.



# Blinded Signatures library

- Available through npm
  - From your project directory, type  
`npm install blind-signatures`
  - Source code/documentation available at  
<https://github.com/kevinejohn/blind-signatures>
- Signatures happen on the *hashes* of the documents

# JavaScript tip of the day: destructuring assignment

- Uses pattern matching to break apart more complex values.
- Useful for returning multiple values.
- (Recent versions of JS only).

# Returning Multiple Values (old school)

```
function foo() {  
    return [1, "one"];  
}
```

```
let arr = foo();  
let n = arr[0];  
let s = arr[1];  
console.log(n); // Prints '1'  
console.log(s); // Prints 'one'
```

# Destructuring Assignment with Arrays

```
function foo() {  
    return [1, "one"];  
}
```

```
let [ n, s ] = foo();
```

```
console.log(n); // Prints '1'  
console.log(s); // Prints 'one'
```

# Destructuring Function Parameters

```
function bar({x, y}) {  
    return x + y;  
}
```

```
let o = {x: 3, y: 4};  
let z = bar(o);
```

```
console.log(z);
```

# Back to Blind Signatures Library ...

- `blind`: returns
  - Blinded hash of the document
  - Its blinding factor
- `sign`: signs blinded document
- `unblind`: removes blinding factor from signature
- `verify`: determines validity of signature
- `verify2`: same as `verify`
  - uses private key (for efficiency I guess?)
- `messageToHash`: calculates hash of message

# Blinding a Document

```
let { blinded, r } =  
  blindSignatures.blind({  
    message: document,  
    N: pubKey.N,  
    E: pubKey.E,  
  });
```

# Signing a Document

```
let signed =  
  blindSignatures.sign({  
    blinded: blinded,  
    key: key,  
  });
```



# Unblinding a Document

```
let unblinded =  
  blindSignatures.unblind({  
    signed: signed,  
    N: pubKey.N,  
    r: r,  
  })
```

# Unblinding a Document

```
let result =  
    blindSignatures.verify({  
        unblinded: unblinded,  
        N: pubKey.N,  
        E: pubKey.E,  
        message: document,  
    });
```

## Lab, Part 2

Details in Canvas and on course website.

Starter code is available on the course website.