# Secret Sharing, Random Numbers, and Information Hiding 

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- Summer CS program held in Yverdon-les-Bains, Switzerland.
- Applications are due April 1.

- More details at http://www.cs.sjsu.edu/su/su14/index.htm



## Secret Sharing: Motivation

- Goal: make secret available, but make it hard to peek.
- Divide secret among multiple organizations.
- Separately, the pieces of give no information about the secret.



# Suppose you want to share a secret number S between Alice and Bob. 

How can you divide it
between them?

## Shamir's Secret Sharing


$\square$ Two points determine a line
$\square$ Give $\left(\mathrm{X}_{0}, \mathrm{Y}_{0}\right)$ to Alice
$\square$ Give ( $\mathrm{X}_{1}, \mathrm{Y}_{1}$ ) to Bob
$\square$ The secret is $(0, S)$, i.e. where the line crosses the y axis.

2 out of 2

## Shamir's Secret Sharing



## Shamir's Secret Sharing



## Shamir's Secret Sharing



2 out of 3
$\square$ Give $\left(\mathrm{X}_{0}, \mathrm{Y}_{0}\right)$ to Alice
$\square$ Give $\left(\mathrm{X}_{1}, \mathrm{Y}_{1}\right)$ to Bob
$\square$ Give ( $\mathrm{X}_{2}, \mathrm{Y}_{2}$ ) to Charlie
$\square$ Then any two can cooperate to find secret $S$
$\square$ But one can't find secret $S$

- A "2 out of 3" scheme


## Shamir's Secret Sharing



3 out of 3
$\square$ Give $\left(\mathrm{X}_{0}, \mathrm{Y}_{0}\right)$ to Alice
$\square$ Give $\left(X_{1}, Y_{1}\right)$ to Bob
$\square$ Give $\left(\mathrm{X}_{2}, \mathrm{Y}_{2}\right)$ to Charlie
$\square 3$ pts determine parabola
$\square$ Alice, Bob, and Charlie must cooperate to find $S$

- A "3 out of 3" scheme
$\square$ What about " 3 out of 4 "?


## Secret Sharing Example

- Key escrow - suppose it's required that your key be stored somewhere
- Key can be "recovered" with court order
- But you don't trust FBI to store your keys
- We can use secret sharing
- Say, three different government agencies
- Two must cooperate to recover the key


## Secret Sharing Example



- Your symmetric key is K
$\square$ Point $\left(\mathrm{X}_{0}, \mathrm{Y}_{0}\right)$ to FBI
$\square$ Point $\left(X_{1}, Y_{1}\right)$ to DoJ
$\square$ Point $\left(X_{2}, Y_{2}\right)$ to DoC
$\square$ To recover your key K, two of the three agencies must cooperate
$\square$ No one agency can get K


## Visual Cryptography

- Another form of secret sharing...
- Alice and Bob "share" an image
- Both must cooperate to reveal the image
- Nobody can learn anything about image from Alice's share or Bob's share
- That is, both shares are required
- Is this possible?


## Visual Cryptography

- How to share a pixel?
- Suppose image is black and white
- Then each pixel is either black or white
- We split pixels as shown



## Visual Cryptography

- How to share a pixel?
- Suppose image is black and white
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## Sharing a B\&W Image

- If pixel is white, randomly choose $a$ or $b$ for Alice's/Bob's shares
- If pixel is black, randomly choose c or d
- No information in one "share"



## Visual Crypto Example

$\square$ Alice's share

$\square$ Bob's share

$\square$ Overlaid shares


## Visual Crypto

- How does visual "crypto" compare to regular crypto?
- In visual crypto, no key...
- Or, maybe both images are the key?
- With encryption, exhaustive search
- Except for a one-time pad
- Exhaustive search on visual crypto?
- No exhaustive search is possible!


## Visual Crypto

- Visual crypto - no exhaustive search...
- How does visual crypto compare to crypto?
- Visual crypto is "information theoretically" secure - true of other secret sharing schemes
- With regular encryption, goal is to make cryptanalysis computationally infeasible
- Visual crypto an example of secret sharing
- Not really a form of crypto, in the usual sense


## Could we design a secret sharing system using one-time pads?



# Random Numbers in Cryptography 

## 808

## "Random" Numbers

- Random numbers are widely used outside of security:
- statistical modeling
- simulations
- random samplings
- For these uses, numbers need to be "statistically random" (they need to appear to be random).


## Random Numbers in Security

- Random numbers used to generate keys
- Symmetric keys
- RSA: Prime numbers
- Diffie Hellman: secret values
- Random numbers used for nonces
- Sometimes a sequence is OK
- But sometimes nonces must be random
- These numbers must be difficult to guess.


## Random Numbers

- Cryptographic random numbers must be statistically random and unpredictable
- Suppose server generates symmetric keys...
- Alice: $\mathrm{K}_{\mathrm{A}}$
- Bob: $\mathrm{K}_{\mathrm{B}}$
- Charlie: $\mathrm{K}_{\mathrm{C}}$
- Dave: $\mathrm{K}_{\mathrm{D}}$
- But, Alice, Bob, and Charlie don't like Dave
- Alice, Bob, and Charlie working together must not be able to determine $\mathrm{K}_{\mathrm{D}}$


## Non-random Random Numbers

- Online version of Texas Hold 'em Poker
o ASF Software, Inc.

- Random numbers used to shuffle the deck
- Program did not produce a random shuffle
- A serious problem or not?


## Card Shuffle

- There are $52!>2^{225}$ possible shuffles
- The poker program used "random" 32-bit integer to determine the shuffle
- So, only $2^{32}$ distinct shuffles could occur
- Code used Pascal pseudo-random number generator (PRNG): Randomize()
- Seed value for PRNG was function of number of milliseconds since midnight
- Less than $2^{27}$ milliseconds in a day
- So, less than $2^{27}$ possible shuffles


## Card Shuffle

- Seed based on milliseconds since midnight
- PRNG re-seeded with each shuffle
- By synchronizing clock with server, number of shuffles that need to be tested $<2^{18}$
- Could then test all $2^{18}$ in real time
- Test each possible shuffle against "up" cards
- Attacker knows every card after the first of five rounds of betting!


## Poker Example

- Poker program is an extreme example
- But common PRNGs are predictable
- Only a question of how many outputs must be observed before determining the sequence
- Crypto random sequences not predictable
- For example, keystream from RC4 cipher
- But "seed" (or key) selection is still an issue!
- How to generate initial random values?
- Keys (and, in some cases, seed values)


## What is Random?

- True "randomness" hard to define
- Entropy is a measure of randomness
- Good sources of "true" randomness
- Radioactive decay - radioactive computers are not too popular
- Hardware devices - many good ones on the market
- Lava lamp - relies on chaotic behavior


## Randomness

- Sources of randomness via software
- Software is (hopefully) deterministic
- So must rely on external "random" events
- Mouse movements, keyboard dynamics, network activity, etc., etc.
- Can get quality random bits by such methods
- But quantity of bits is very limited


## The Bottom Line

"The use of pseudo-random
processes to generate secret quantities can result in pseudosecurity"

## Information Hiding

A boat, beneath a sunny sky Lingering onward dreamily In an evening of July Children three that nestle near, Eager eye and willing ear, — Lewis Carroll, Through the Looking Glass

A boat, beneath a sunny sky Lingering onward dreamily

## In an evening of July -

 Children three that nestle near, Eager eye and willing ear,— Lewis Carroll, Through the Looking Glass

## Information Hiding

- Digital Watermarks
- Example: Add "invisible" identifier to data
- Defense against music or software piracy
- Steganography
- "Secret" communication channel
- Similar to a covert channel (more on this later)
- Example: Hide data in image or music file


## Watermark

- Add a "mark" to data
- Visibility of watermarks
- Invisible - Watermark is not obvious
- Visible - Such as TOP SECRET
- Robustness of watermarks
- Robust - Readable even if attacked
- Fragile - Damaged if attacked


## Watermark Examples

- Add robust invisible mark to digital music
- If pirated music appears on Internet, can trace it back to original source of the leak
- Add fragile invisible mark to audio file
- If watermark is unreadable, recipient knows that audio has been tampered (integrity)
- Combinations of several types are sometimes used
- E.g., visible plus robust invisible watermarks


## Watermark Example (1)

- Non-digital watermark: U.S. currency

$\square$ Image embedded in paper on rhs
o Hold bill to light to see embedded info


## Watermark Example (2)

- Add invisible watermark to photo
- Claimed that 1 inch $^{2}$ contains enough info to reconstruct entire photo
- If photo is damaged, watermark can be used to reconstruct it!


## Steganography

- According to Herodotus (Greece 440 BC )
- Shaved slave's head
- Wrote message on head
- Let hair grow back
- Send slave to deliver message
- Shave slave's head to expose message - warning of Persian invasion
- Historically, steganography used more often than cryptography


## Images and Steganography

- Images use 24 bits for color: RGB
- 8 bits for red, 8 for green, 8 for blue
- For example
$-0 \times 7 \mathrm{E} 0 \times 520 \times 90$ is this color
- 0xFE 0x52 0x90 is this color
- While
- 0xAB 0x33 0xF0 is this color
- 0xAB 0x33 0xF1 is this color
- Low-order bits don't matter...


## Images and Stego

- Given an uncompressed image file...
- For example, BMP format
- ...we can insert information into low-order RGB bits
- Since low-order RGB bits don't matter, result will be "invisible" to human eye
- But, computer program can "see" the bits


## Stego Example 1



- Left side: plain Alice image
- Right side: Alice with entire Alice in Wonderland (pdf) "hidden" in the image


## Non-Stego Example

## $\square$ Walrus.html in web browser

"The time has come," the Walrus said, "To talk of many things:
Of shoes and ships and sealing wax Of cabbages and kings And why the sea is boiling hot And whether pigs have wings."

- "View source" reveals:
<font color=\#000000>"The time has come," the Walrus said,,</font><br> <font color=\#000000>"To talk of many things: </font><br> <font color=\#000000>0f shoes and ships and sealing wax </font><br> <font color=\#000000>0f cabbages and kings </font><br>
<font color=\#000000>And why the sea is boiling hot </font><br>
<font color=\#000000>And whether pigs have wings." </font><br>


## Stego Example 2

$\square$ stegoWalrus.html in web browser

> "The time has come," the Walrus said, "To talk of many things:
> Of shoes and ships and sealing wax Of cabbages and kings And why the sea is boiling hot And whether pigs have wings."

- "View source" reveals:
<font color=\#000101>"The time has come," the Walrus said,</font><br>
<font color=\#000100>"To talk of many things: </font><br>
<font color=\#010000>0f shoes and ships and sealing wax </font><br>
<font color=\#010000>0f cabbages and kings </font><br>
<font color=\#000000>And why the sea is boiling hot </font><br>
<font color=\#010001>And whether pigs have wings." </font><br>
$\square$ "Hidden" message: 011010100100000101


## Using Steganography



## Some formats (e.g. HTML) are easier for people to read, but no harder for computers.



We can hide info in unimportant bits, but Trudy can overwrite those bits as well.

## Steganography

- Easy to hide info in unimportant bits
- Easy to destroy info in unimportant bits
- To be robust, must use important bits
- But stored info must not damage data
- Collusion attacks are another concern
- Robust steganography is tricky!


## Information Hiding: The Bottom Line

- Not-so-easy to hide digital information
- "Obvious" approach is not robust
- Stirmark: tool to make most watermarks in images unreadable without damaging the image
- Stego/watermarking active research topics
- If information hiding is suspected
- Attacker may be able to make information/watermark unreadable
- Attacker may be able to read the information, given the original document (image, audio, etc.)

