

# Arithmetic Circuits

CS255

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Feb. 22, 2006.

# Outline

- Introduction
- Combination Circuits
- Full Adders
- Circuits for Addition (start)

# Introduction

- We are now going to consider our second model of parallel computations: circuits.
- The algorithms we are going to demonstrate can be parallelized, are ones for addition, multiplication, etc.
- One motivation in trying to make these parallel is for speeding up the implementation circuits of these operations in hardware.
- Running time here will be depth of the circuit.

# Combinational Elements

- Our circuits will be built out of combinational elements interconnected by wires.
- A **combinational element** is any circuit element that has a constant number of inputs and outputs and that computes a well defined function.
- If the inputs and outputs are boolean ( $\{0,1\}$ ) then we will call an element a **boolean combinational element**.
- Example, could have  $ADD(x,y)$  whose inputs are single bits and whose output is two wire which has there sum in binary. Could view the high order and low order outputs as separate functions.
- If our circuit computes just one boolean function, we call it a **logic gate**. For example, AND, OR, NOT, XOR gates.

# Truth table of a Logic Gate

- Is a table which says what the logic gate does on each possible input:

X	Y	AND
0	0	0
0	1	0
1	0	0
1	1	1

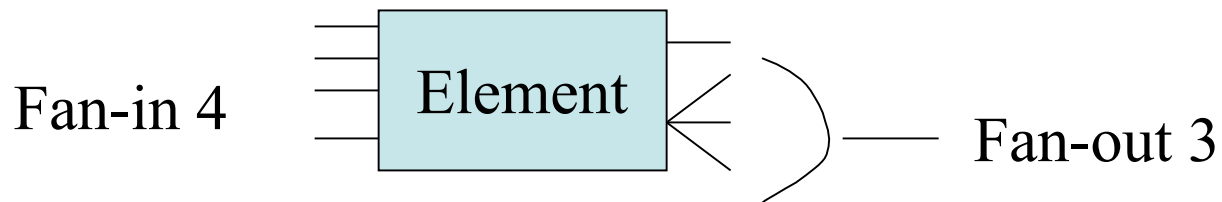
- How many logic possible n input boolean functions?

# Time to Compute a Logic Gate

- Real logic gates take some time to settle and become stable on their output values after inputs values have been applied.
- This fixed amount of time we will call the **propagation delay** of the element.
- We will assume it is constant.

# Combinational Circuits

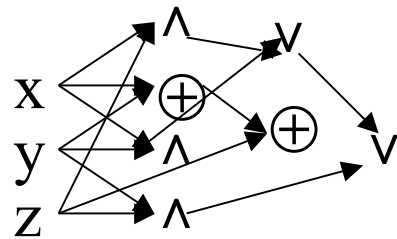
- A combinational circuit will consist of one or more combinational elements interconnected in an acyclic fashion.
- Interconnections will be called **wires**.
- A given output of an element can be connected to several other elements. the number of other elements is its **fan-out**:



- Fan-in is the number of wires going into an element.
- A wire which connects does not connect to an element input is called a **circuit output**.

# Full Adders

- Consider the full adder on three boolean inputs  $x$ ,  $y$ ,  $z$ .
- This functions takes the three inputs and outputs two bits which represent the sum of the inputs in binary.
- The low order output could be computed as the parity on the inputs  $s = \text{parity}(x, y, z) = x \oplus y \oplus z$ .
- The high order output is on if the majority of the inputs are on:  $c = \text{majority}(x, y, z) = (x \wedge y) \vee (y \wedge z) \vee (x \wedge z)$ .
- Here is what a circuit might look like:



The time to compute this function is the length of the longest path = 3.



# Size and Depth

- The **depth** of an input wire is 0.
- The **depth** of an element is the maximum depth of its inputs wire + 1.
- The **depth** of an output wire is the depth of the element it come from.
- The **size** of a circuit is the number of elements it contains.
- So in the last slide the depth of the full adder was 3 and its size was 7.
- Elements of the same depth can operate in parallel.

# Addition Circuits

- We now try to develop efficient circuits for addition of two  $n$  bit numbers.
- For example, we will first look at ripple carry addition which can add two  $n$ -bit numbers in  $\Theta(n)$  size and depth.
- This kind of addition thus does not give us any parallelism.
- We will then look at carry-lookahead addition which also has  $\Theta(n)$  size but not  $O(\log n)$  depth. So very parallel.