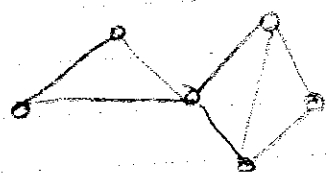


## I Broadcast Links:

A single communication is shared by all machines on the network. Short messages called packets are sent by any machine and are received by all others (called multicast if subset of all others). An address on the packet says who it was intended for.

## Point to Point Links:

These consist of many connections between pairs of machines. A packet might have to traverse several edges to get from its source to target.



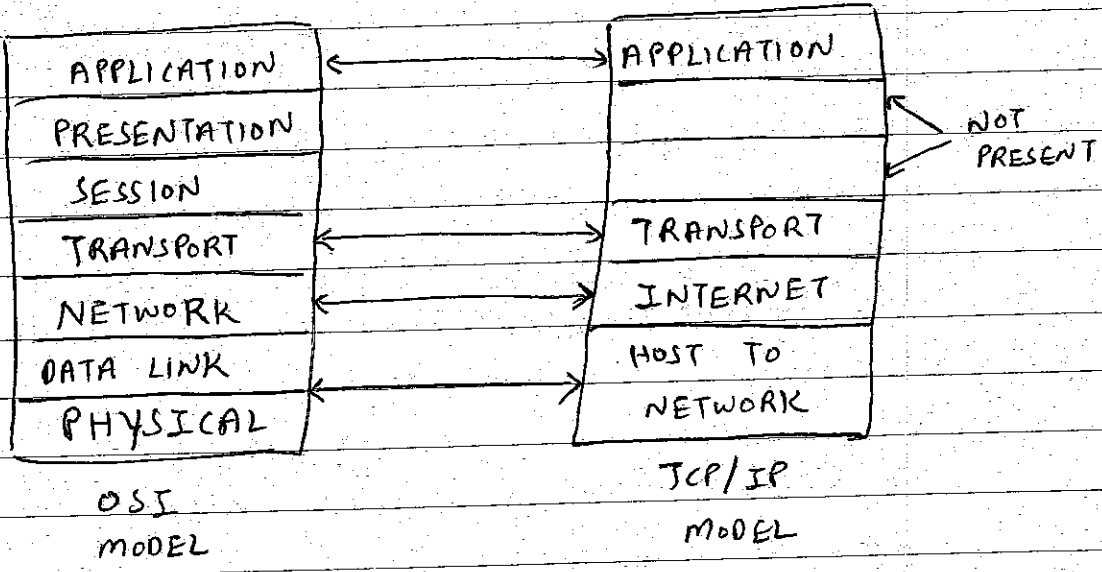
## II

Static - typically divide time into discrete intervals and use round-robin algorithm allowing each machine only to broadcast when its time slot comes up.

### Dynamic:

Centralized - Allocation is a single bus arbitrator decides who gets to go next.

Decentralized: Each machine must decide whether or not to transmit.

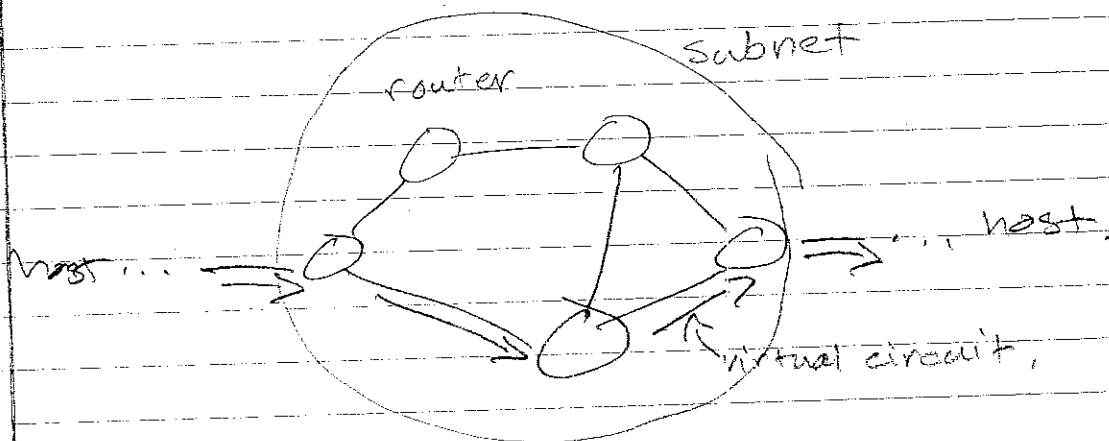


4

cell

5 byte header } Frame.  
48 byte data }

- 1) sends first cell to est. connection, id
- 2) each router est a ~~map~~ cell map upon first cell sent
- 3) creates a virtual circuit.



Notes

- no sequencing needed
- 4 layer 3D reference model

5

According to Nyquist theorem, a noiseless channel with a bandwidth of  $H$  can transmit a maximum data rate of

$$2H \log_2 V$$

Where

$V$  is discrete levels

Given

$$\text{Bandwidth } (H) = 5000 \text{ Hz}$$

$$\text{Discrete levels } (V) = 128$$

$$\rightarrow 2 \cdot 5000 \log_2 128$$

$$= 10,000 \log_2 2^7$$

$$= 70,000 \text{ Bits/sec.}$$

5. Bandwidth = 5 kHz  
up to 128 discrete levels

$$2^5 = 32$$
$$2^6 = 64$$
$$2^7 = 128$$

$$\begin{aligned} \text{max data rate} &= 2B \log_2 V \\ &= 2(5000) \log_2 (128) \\ &= (10\,000)(7) \\ &= 70\,000 \text{ bps} \\ &= 70 \text{ kbps} \end{aligned}$$

6. An active repeater in a fiber network consists of an optical receiver, some copper wires, a signal regenerator, and an optical transmitter. The light from an optical fiber is converted into an electrical signal by the optical receiver, then transmitted through a copper wire to an electrical signal regenerator, where it is regenerated to full strength & resent through copper to an optical transmitter. out the other way.

(7)

Frequency Division Multiplexing:

Frequencies of different channels are translated by different amounts to put them all on the same line.

Time Division Multiplexing

In this scheme, each user takes turns (Round-Robin), each one getting bandwidth for a small period of time.

Wavelength Division Multiplexing

Input signals at different wavelengths are linearly added together by a combiner. At other end, split out by a splitter.

Name: Anand Sivaramakrishnan

Q8) A's chip sequence: 01010101  
B's chip sequence: 11110000

Bipolar Chip Sequence for:

$$\begin{array}{l} A = -1 +1 -1 +1 -1 +1 -1 +1 \\ B = +1 +1 +1 +1 -1 -1 -1 -1 \end{array}$$

~~To see:~~

A wants to send a '1' and  
B wants to send a '0'

∴ A will send its original sequence &  
B will send its inverted sequence.

$$\begin{array}{l} \text{For A: } -1 +1 -1 +1 -1 +1 -1 +1 \\ \text{for B: } -1 -1 -1 -1 +1 +1 +1 +1 \\ \text{Addition} \quad \underline{-2 \quad 0 \quad -2 \quad 0 \quad 0 \quad 2 \quad 0 \quad 2} \end{array}$$

Signal to be  
Transmitted.

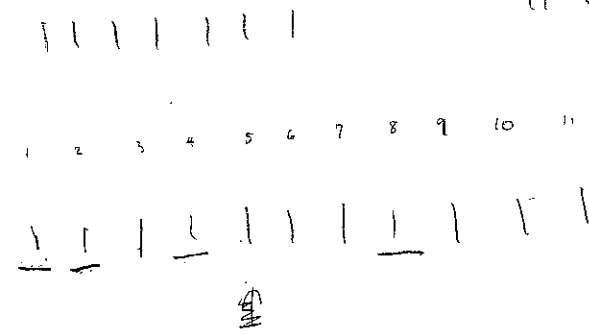
At receiver,

To decode A, multiply the signal  
received with the bipolar chip  
sequence of 'A'

$$\begin{aligned} & (-2 \quad 0 \quad -2 \quad 0 \quad 0 \quad 2 \quad 0 \quad 2) \begin{pmatrix} -1 +1 -1 +1 \\ -1 +1 -1 +1 \end{pmatrix} \\ & = (2 \quad 0 \quad 2 \quad 0 \quad 0 \quad 2 \quad 0 \quad 2) \\ & = 8, \text{ and then divide by } 8 \\ & 8/8 = 1 // \text{ received for A} \end{aligned}$$

9)

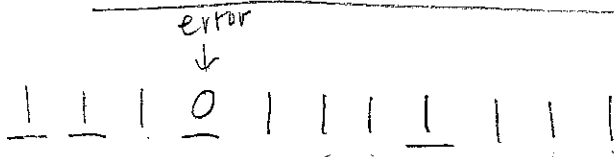
11 bit Hamming



- 3 = 0011
- ~~4 = 0100~~
- 5 = 0101
- 6 = 0110
- 7 = 0111
- 9 = 1001
- 10 = 1010
- 11 = 1011

0011  
0101  
1001  
1011

count:  $3355$   
( $2^3 2^2 2^1 2^0$ )



~~position 4 is parity bit which should be 1~~

~~so follow~~

~~position 4 should be 1~~

parity  
Calculated  
from data  
bits

$2^3 \quad 2^2 \quad 2^1 \quad 2^0$

| | | |

parity  
in codeword

| 0 | | |

value:

0 1 0 0  $\Rightarrow$   $2^4$   $\Rightarrow$  position 4 is incorrect.





10

A

B

A sends  $(0, 1, A0)$

→ B gets  $(0, 1, A0)^*$

B sends  $(0, 0, B0)$

A gets  $(0, 0, B0)^*$

B sends  $(0, 0, B0)$   
(TIMEOUT)

A sends  $(1, 0, A1)$

(A rejects  $B0$ ) ←

→ B gets  $(1, 0, A1)^*$

- Machine A is slower than Machine B  
So,  $delay_A > timeout_B$