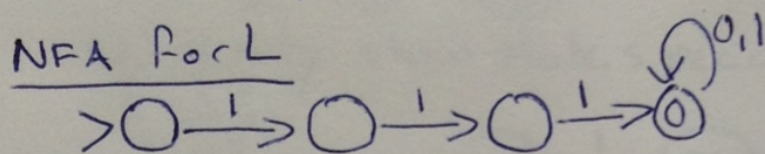


HW 2 Grp 1

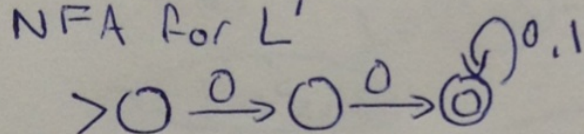
- ① Let $L = \{w \mid w \text{ begins with three 1's}\}$
 $L' = \{w \mid w \text{ begins with two 0's}\}$.

Draw NFAs for both languages. Then use the constructions from class to make NFAs for LL' , $L \cup L'$ and L^* .

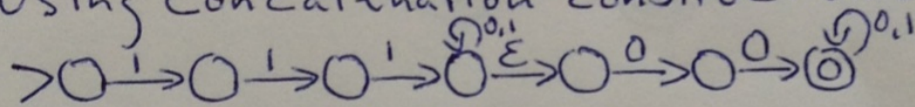
NFA for L



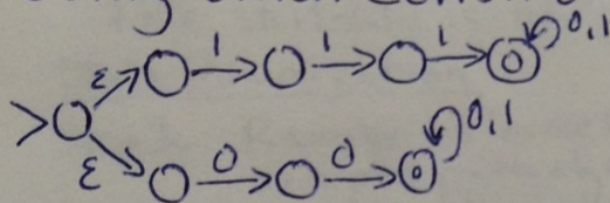
NFA for L'



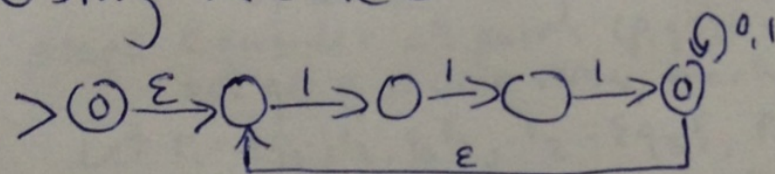
Using concatenation construction:



Using union construction:

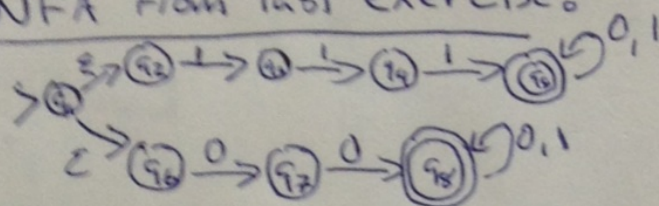


Using Kleene Star construction:

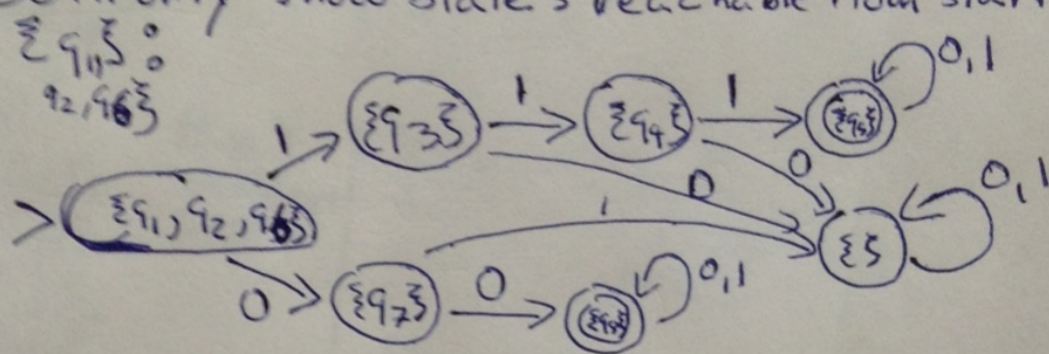


2 Use the power set construction to convert your NFA from the last exercise for $L \cup L'$ into a DFA then do minimization on the result.

NFA from last exercise:



Will only show states reachable from start state



Note start state is $E(\{q_1\}) = \{q_1, q_2, q_6\}$.
State Minimization

Step 1 Remove all inaccessible states using simple path algorithm.
 we already implicitly did this above by not considering all states in power set.

Step 2 Consider all pairs (p, q) of states. If $p \in F$ but $q \notin F$ or vice versa, mark (p, q) distinguishable

Let $r_1 = \{q_1, q_2, q_6\}$, $r_2 = \{q_3\}$, $r_3 = \{q_4\}$, $r_4 = \{q_5\}$,
 $r_5 = \{q_7\}$, $r_6 = \{q_8\}$, $r_7 = \{\epsilon\}$, so have simpler names for states. Then get:

	r_1	r_2	r_3	r_4	r_5	r_6	r_7
r_1	X			X		X	
r_2		X		X		X	
r_3			X	X		X	
r_4	X	X	X	X	X	X	X
r_5				X	X	X	
r_6	X	X	X		X	X	X
r_7				X		X	X

Cprob2 cont'd)

next your

Step 3 Repeat until no previously unmarked pairs are marked

(a) For all (p, q) and all $a \in \Sigma$, compute $\delta(p, a) = p_a$ and $\delta(q, a) = q_a$. If (p_a, q_a) is marked as distinguishable, mark (p, q) as distinguishable

Pass 1 applying (a)

	r_1	r_2	r_3	r_4	r_5	r_6	r_7
r_1			X	X	X	X	X
r_2			X	X	X	X	X
r_3	X	X		X	X	X	X
r_4	X	X			X		X
r_5	X	X	X	X		X	X
r_6	X	X	X		X		X
r_7	X	X	X	X	X	X	

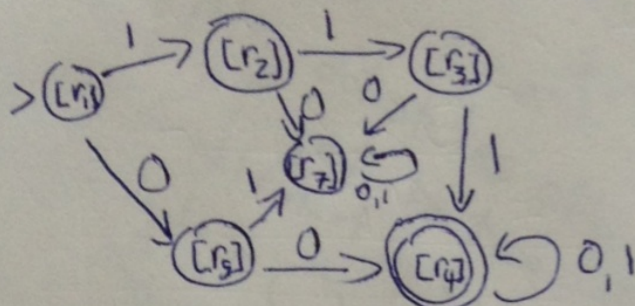
Example why (r_5, r_7) marked this round
because $\delta(r_5, 0) = r_6$ (accepting)
 $\delta(r_7, 0) = r_6$ (accepting)
 (r_6, r_7) marked previously

Pass 2 applying (a)

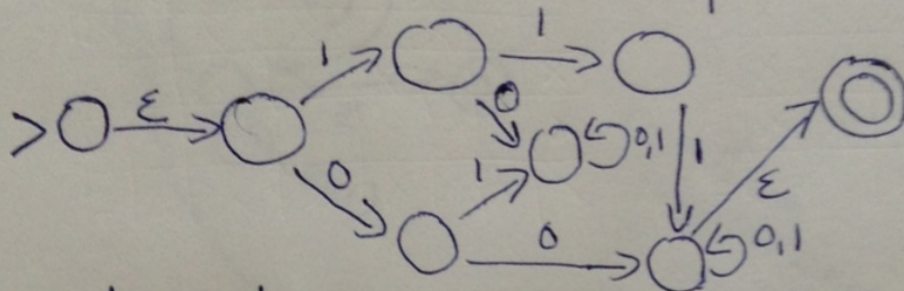
	r_1	r_2	r_3	r_4	r_5	r_6	r_7
r_1		X	X	X	X	X	X
r_2	X		X	X	X	X	X
r_3	X	X		X	X	X	X
r_4	X	X	X		X		X
r_5	X	X	X	X		X	X
r_6	X	X	X		X		X
r_7	X	X	X	X	X	X	

Example why (r_1, r_2) marked
 $\delta(r_1, 1) = r_2$
 $\delta(r_2, 1) = r_3</$

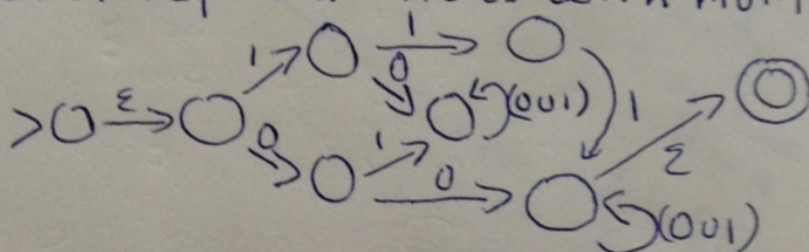
- ③ Use the construction from class to convert your DFA to a regular expression.
DFA from last exercise



1st add new start and accept states



2nd replace arrows with multiple labels with union



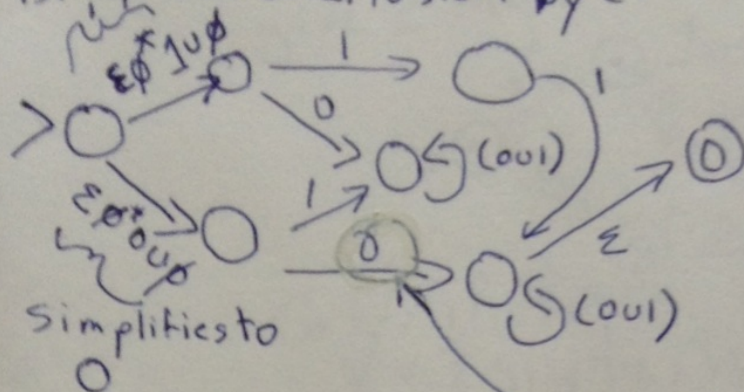
3rd add $\emptyset$ labels as per construction



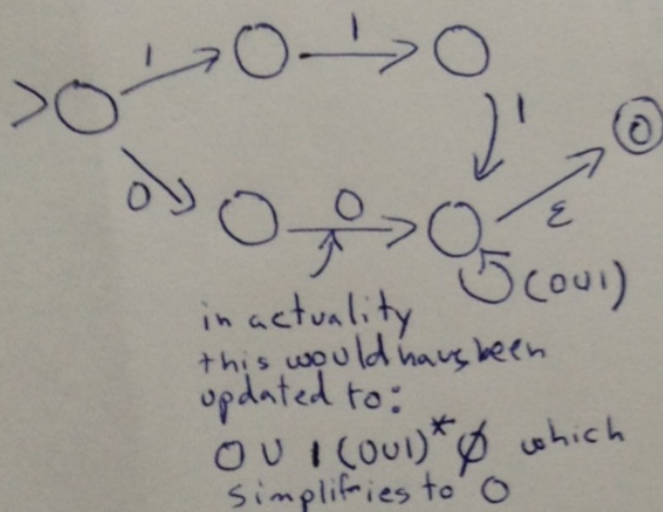
← very messy
 will suppress $\emptyset$
 states on
 next page
 so can
 read
 diagrams

3 cont'd

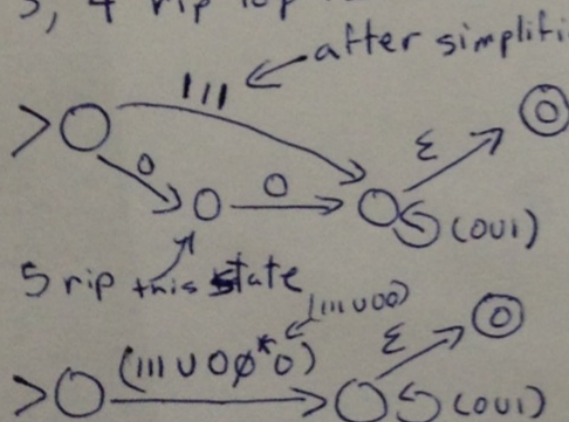
Start Ripping out states other than first and last till get to two states
1st rip state connect to start by ϵ .



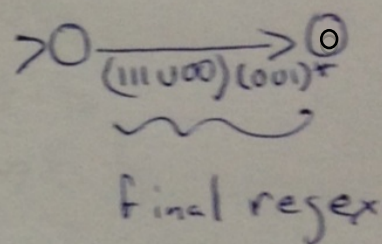
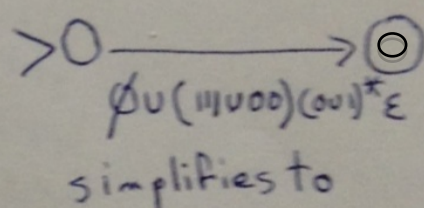
2nd rip trapstate



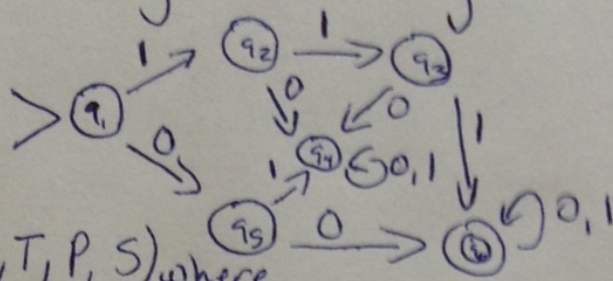
3, 4 rip top two 1 transition states



Rip middle state



4 Using the construction from class
give a right-linear grammar for the DFA:



$G = (V, T, P, S)$ where

$$V = \{q_1, q_2, q_3, q_4, q_5, q_6\}$$

$$T = \{0, 1\}$$

$$P = \left\{ \begin{array}{l} q_1 \rightarrow 1q_2, \\ q_1 \rightarrow 0q_5, \\ q_2 \rightarrow 1q_3, \\ q_2 \rightarrow 0q_4, \\ q_3 \rightarrow 1q_6, \\ q_3 \rightarrow 0q_4, \\ q_4 \rightarrow 1q_4, \\ q_4 \rightarrow 0q_4, \\ q_5 \rightarrow 1q_4, \\ q_5 \rightarrow 0q_6, \\ q_6 \rightarrow 0q_6, \\ q_6 \rightarrow 1q_6, \\ q_6 \rightarrow \epsilon \end{array} \right\}$$

$$S = q_1$$